

Entropy Loss and Random Functions

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Outline

1 Introduction

2 Entropy Loss

3 Attacks

4 Open Problems

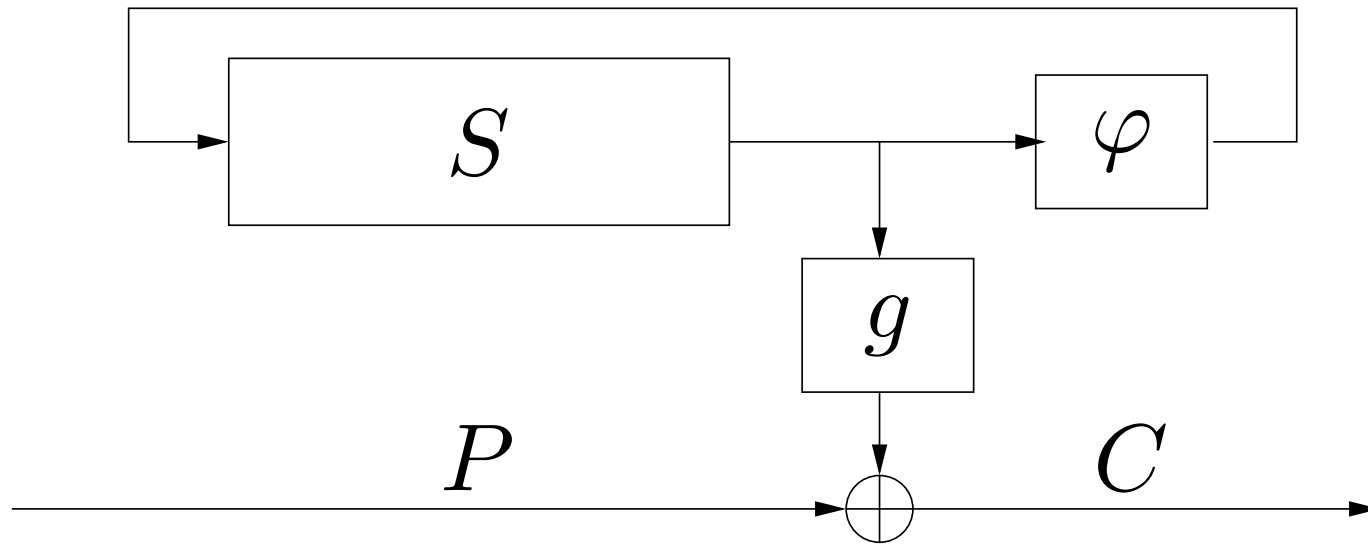
Part 1

Introduction

Random Function

By *random function* we refer to a function φ which was randomly chosen out of $\mathcal{F}_n = \{\varphi : \Omega_n \rightarrow \Omega_n\}$, where Ω_n is a set of n elements.

Stream Cipher Model



- S state with $s_t \in \Omega_n, t \geq 0$
- $\varphi \in \mathcal{F}_n$ random function which updates S
- P plaintext
- C ciphertext
- g filter function

Entropy of the State

- $\{p_i\}_{i=1}^n$ distribution of initial state
- $p_i^\varphi(t)$ probability of i after t iterations of φ
- entropy:

$$H_t^\varphi = \sum_{i=1}^n p_i^\varphi(t) \log_2 \left(\frac{1}{p_i^\varphi(t)} \right)$$

- $E(H_t)$ expected entropy after k iterations

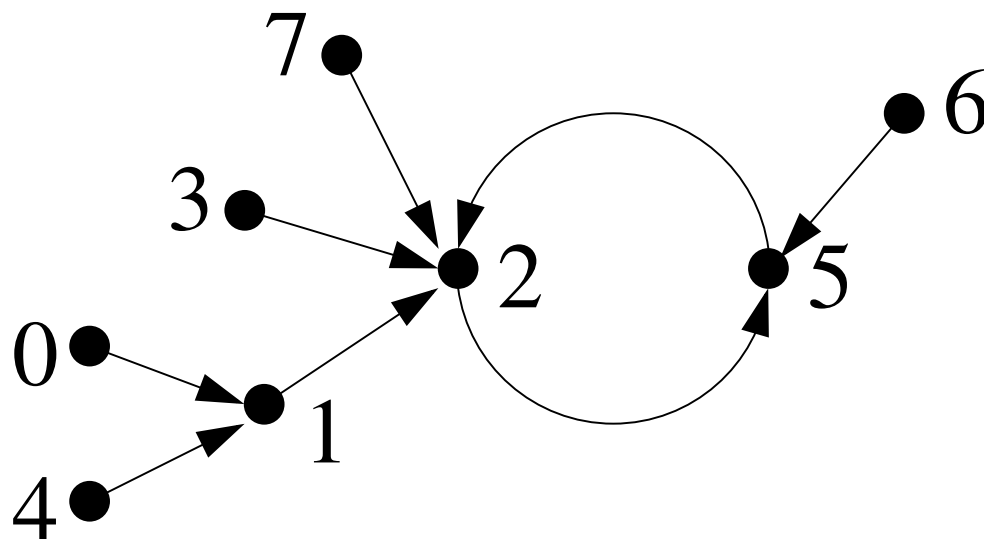
Part 2

Entropy Loss

Properties of Random Functions (1)

[Flajolet, Odlyzko 98]

$$f(x) = x^2 \pmod{8}$$



Properties of Random Functions (2)

[Flajolet, Odlyzko 98]

- Expected **number of image points** after k iterations is

$$(1 - \tau_k)n$$

with $\tau_0 = 0$ and $\tau_{k+1} = e^{\tau_k - 1}$.

- Expected **cicle points** $cp(n) = \sqrt{\frac{\pi n}{2}}$.
- Expected **maximal tail length**
 $mtl(n) = \sqrt{n} 2\pi \log(2)$.

Properties of Random Functions (3)

[Flajolet, Odlyzko 98]

- Expected number of r -nodes after one iteration is

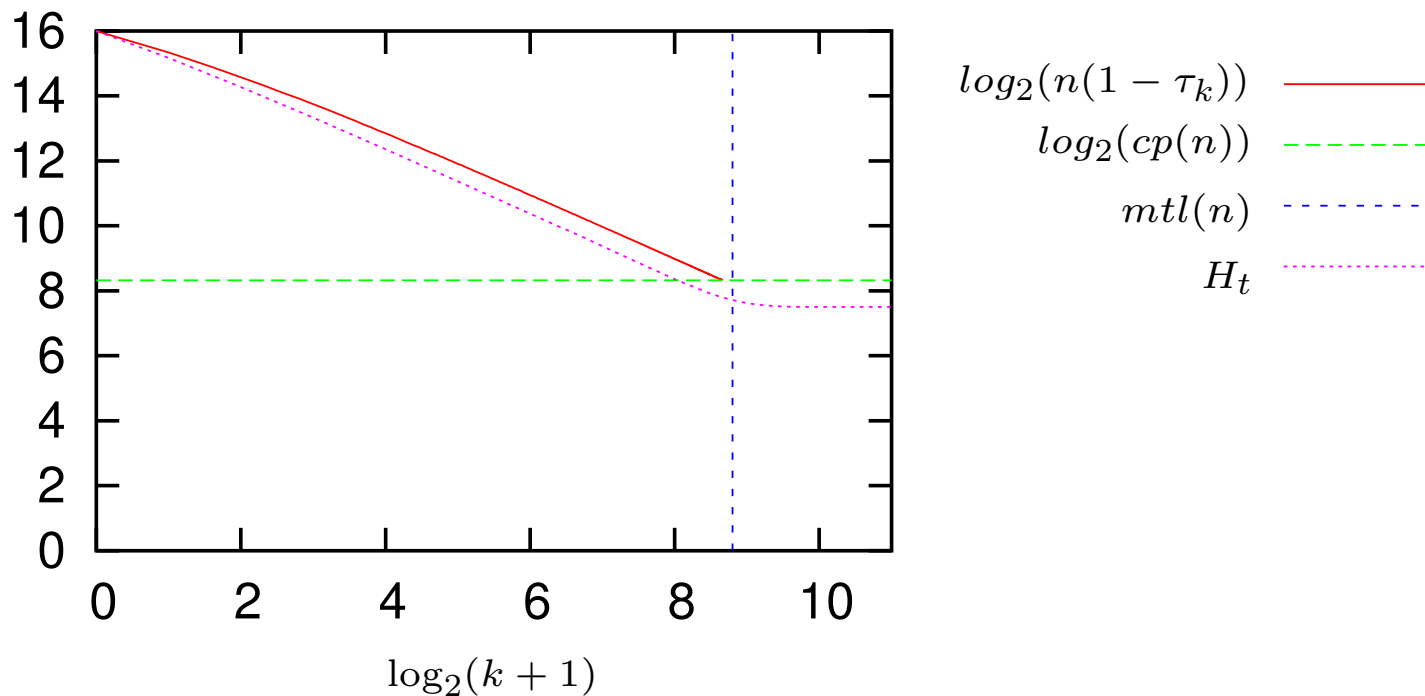
$$\frac{n}{e r!}$$

(By an r -node we refer to a node with r incoming edges in the functional graph or respectively a preimage of size r)

Upper bound with image points

$$E(H_t) \leq \log_2(n) - \log_2(1 - \tau_k)$$

- for
- $k \leq mtl(n)$ and
 - $(1 - \tau_k)n \geq cp(n)$.



Entropy by means of r nodes

- **Uniform** initial distribution

$$E(H_1^U) = \frac{n}{e} \sum_{r=1}^n \frac{1}{r!} \frac{r}{n} \log_2 \frac{r}{n}$$

- **Arbitrary** initial distribution

$$\frac{1}{\binom{n}{r}} \sum_{1 \leq i_1 < \dots < i_r \leq n} (p_{i_1} + \dots + p_{i_r}) \log_2 \frac{1}{p_{i_1} + \dots + p_{i_r}}$$

Comparison for $n = 2^{16}$

Method	Approximation	$n = 2^{16}$
image points	$\log(n) - 0.6617$	15.3383
r nodes (uniform)	$\log(n) - 0.8272$	15.1728
test		15.1728

Part 3

Attacks