# Recovering a code's length and synchronization from a noisy intercepted bitstream. 

M. Cluzeau and M. Finiasz

## Overview of the problem


$\approx$ We intercept a noisy bitstream and want to recover the (encrypted) information.

## Overview of the problem

- Most of the time, coding schemes are standardized
$\triangleright$ no need for code reconstruction.
- Yet, "some people" are interested in this:
$\triangleright$ not many public works on this topic,
$\triangleright$ many interesting problems arise, depending on the type of code we focus on.
- Here we focus on linear block codes requiring to:
$\triangleright$ find the block length,
$\triangleright$ find a generator/parity check matrix,
$\triangleright$ find an efficient decoder,
$\rightarrow$ we do not address this problem here.


# Overview of the problem 

The case of linear block codes

## $S$ <br> $010010110100001011101010010100101110 .$.

- The only thing we have is a noisy bitstream:
$\triangleright$ we need to find $s_{0}$ and $n_{0}$ the synchronization and length of the code.
- For very short codes of small dimension various techniques can give us some hint on $n$,
$\ddagger \quad \triangleright$ none of them work for real life codes...
$\rightarrow$ we have to test each choice of $s$ and $n$.


## In the absence of noise

- For given $s$ and $n$ build the matrix $\mathcal{G}$ of "codewords"
$\triangleright$ if $n=n_{0}$ and $s=s_{0}$ it has minimal rank $k$,
$\triangleright$ if $n=n_{0}$ and $s \neq s_{0}$ it has rank $\min \left(k+\left|s-s_{0}\right|, n\right)$,
$\triangleright$ if $n \neq n_{0}$ it has rank $n$.



## In the absence of noise

## The easy case...

- For given $s$ and $n$ build the matrix $\mathcal{G}$ of "codewords"
$\triangleright$ if $n=n_{0}$ and $s=s_{0}$ it has minimal rank $k$,
$\triangleright$ if $n=n_{0}$ and $s \neq s_{0}$ it has rank $\min \left(k+\left|s-s_{0}\right|, n\right)$,
$\triangleright$ if $n \neq n_{0}$ it has rank $n$.
- Very efficient to guess $n_{0}$ and then $s_{0}$, $\rightarrow$ only for very low noise levels $\tau \ll \frac{1}{n}$.
- For higher noises the rank is always $n \ldots$


## In the presence of noise

 Using words of the dual- If $n$ and $s$ are correct, a word of the dual of the target code multiplied by $\mathcal{G}$ should have low weight,
$\triangleright$ suppose we have such a dual word of low weight $w$.



# In the presence of noise 

 Using words of the dual- If a word following the green distribution is found, $n=n_{0}$ $\triangleright$ and $s-s_{0}$ is probably small.



## The algorithm we propose

- We need to exhaustively search through the possible $s$ and $n$.
- Successively go through the possible values of $n$ $\triangleright$ for each length "test" several synchronizations $s$ $\rightarrow$ different possible heuristics.
- Testing a pair $(n, s)$ consists in searching for a dual word following the green distribution:
$\triangleright$ exhaustive search of words of given weight
$\triangleright$ using Valembois' algorithm.


## Exhaustive search of given weight dual words

- We look for a dual word of length $n$ and weight $w$.
- We can find all such dual words using:
$\triangleright$ straight-forward exhaustive search
$\rightarrow O\left(n^{w}\right)$ time and $0(1)$ memory.
$\triangleright$ the birthday algorithm
$\rightarrow O\left(n^{\frac{w}{2}}\right)$ time and $O\left(n^{\frac{w}{2}}\right)$ memory.
$\triangleright$ the Chose-Joux-Mitton algorithm [Eurocrypt 2002]
$\rightarrow O\left(n^{\frac{w}{2}}\right)$ time and $O\left(n^{\left[\frac{w}{4}\right\rceil}\right)$ memory.
® Very efficient for codes with very low weight dual words $\rightarrow$ typically LDPC codes.


## Valembois' algorithm


$\triangleright$ Based on the Canteaut-Chabaud decoding algorithm,
$\triangleright$ does not focus only on low weight dual words,
$\triangleright$ small memory requirements.

- Very efficient for low noise levels,
$\rightarrow$ tolerates higher noise levels for very short codes.
- Codes of rate $\frac{1}{2}$ :
$\triangleright$ no low weight dual words,
$\triangleright$ for our problem: among the difficult cases.
- Dual words found in 10000 iterations of Valembois' algorithm (less than a second).

| $n$ | 0.001 | 0.002 | 0.005 | 0.01 | 0.02 | 0.05 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 32 | 14637 | 27081 | 42570 | 42913 | 19464 | 210 |
| 64 | $\infty$ | $\infty$ | $\infty$ | 1172189 | 6310 | 0 |
| 128 | $\infty$ | $\infty$ | $\infty$ | 2992 | 0 | 0 |
| 256 | $\infty$ | $\infty$ | 0 | 0 | 0 | 0 |

- LDPC codes of rate $\frac{1}{2}$ and weight 6 parity checks, $\triangleright$ find words for lengths up to 10000 with 2GB memory.
- For an LDPC of length 1000 in 50 iterations ( $\sim 2$ min.)

| $\tau$ | words <br> found | expected words <br> per iteration | expected total <br> words found |
| :---: | :---: | :---: | :---: |
| 0.01 | 478 | 41 | 492 |
| 0.02 | 251 | 7.5 | 266 |
| 0.03 | 84 | 1.5 | 70 |
| 0.04 | 15 | 0.33 | 16 |
| 0.05 | 6 | 0.08 | 3.9 |
| 0.06 | 1 | 0.02 | 1.0 |

## Conclusion

- We can find the length/synchronization of a code by using reconstruction techniques,
$\triangleright$ easier for codes with low weight dual words $\rightarrow$ LDPC
$\triangleright$ not very satisfying for random codes.
- For an unknown code, both techniques should be tried $\triangleright$ for very low noise levels, Valembois' algorithm is faster, even for long LDPC codes.
- For other kind of codes:
$\triangleright$ convolutional codes
[Côte,Sendrier - ISIT09]
$\triangleright$ turbocodes $\rightarrow$ we are working on it...

