Quantum Quadratic form reduction

Quantum multiplication by 2^x

Quantum binary quadratic form reduction A $O(n \log n)$ depth circuit and application to lattice reduction

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Quantum Quadratic form reduction

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Human Context

Nicolas David PHD student at Inria Paris Symmetric Cryptography and Quantum Cryptanalysis



Thomas Espitau Postdoc at NTT in Tokyo Lattice Based Crytpography



Akinori Hosoyamada Researcher at NTT in Tokyo Symmetric Cryptography and Quantum Cryptography

Lattice

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- A lattice is a discrete subgroup of **R**ⁿ
- u_1 , u_2 is a reduced basis

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Scientific Context

Lattice

- Lattice based cryptography is based on the difficulty of finding a reduced basis
- LLL is an algorithm that computes a reduced basis

Quantum Cryptanalysis

 With the help of quantum computer it is possible to break some cryptographic primitives

Quantum Quadratic form reduction

Quantum multiplication by 2^{x}

How it all started



Quantum Quadratic form reduction

Quantum multiplication by 2^{x}

How it all started



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Efficient reversible shallow circuits

Designing efficient shallow circuits is a new domain of research. It mainly consists in

- Designing efficient circuit for basic operations
- Give its precise complexity allowing to derive precise costs in applications

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State of the art

Some previous work

- Addition
- Multiplication
- GCD

Today we will look into Binary quadratic form reduction

- first circuit designed for quadratic form reduction
- Application to LLL

⁰Yasuhiro Takahashi and Noboru Kunihiro. "A fast quantum circuit for addition with few qubits. Quantum Information Computation" Srijit Dutta, Debjyoti Bhattacharjee, and Anupam Chattopadhyay. "Quantum circuits for toom-cook multiplication." Mehdi Saeedi and Igor L Markov. "Quantum Circuits for GCD Computation with O(nlogn) Depth and O(n) Ancillae."

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1 Preliminaries

- Quantum Circuits
- Binary quadratic form
- Quantum Quadratic form reduction
- 3 Quantum multiplication by 2^x

Quantum Quadratic form reduction

Quantum multiplication by 2^x

Outline

1 Preliminaries

Quantum Circuits

• Binary quadratic form

Quantum Quadratic form reduction

3 Quantum multiplication by 2[×]

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Quantum circuit

Qubit

A qubit is the basic unit of quantum information

- $\bullet\,$ pure states : $|0\rangle$, $|1\rangle$ (extended to $|00\rangle$, $|01\rangle$, $|11\rangle\,$ etc.)
- superposition : $|\psi\rangle = \alpha |0\rangle + \beta |1\rangle$ with $\alpha, \beta \in \mathbf{C}$ and $|\alpha|^2 + |\beta|^2 = 1$

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Some Quantum Gates

• NOT.

$$|x\rangle \longrightarrow |\bar{x}\rangle$$

• Bitwise addition (CNOT).



• Toffoli (CCNOT).



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Quantum multiplication by 2^x

Some Quantum Gates

• Swap.



Quantum Quadratic form reduction

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Quantum circuit and complexity measures



 $\left|0\right\rangle^{k}$ are called Ancilae qubits

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Quantum circuit and complexity measures



- $\left|0\right\rangle^{k}$ are called Ancilae qubits
 - Depth : Related to the speed of the execution.
 - Width : Related to the memory (# Ancilae).
 - Volume: Total number of operations.

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Outline



- Binary quadratic form
- Quantum Quadratic form reduction
- Quantum multiplication by 2[×]

Quantum Quadratic form reduction

Quantum multiplication by 2^x

Binary quadratic form : Definition

Definition

An (integral) Binary quadratic form Q = [A, B, C] is a polynomial $(AX^2 + BXY + CY^2) \in [X, Y]$. The integer $\Delta = B^2 - 4AC$ is called the *discriminant* of Q. The form is said to be:

- Degenerate when $\Delta = 0$
- Positive (resp. Negative) Definite when $\Delta < 0$ and $Q(x, y) \ge 0$ (resp. $Q(x, y) \le 0$) for any $(x, y) \in \mathbf{R}^2$.

Associated matrix :
$$Q = \begin{pmatrix} A & B/2 \\ B/2 & C \end{pmatrix}$$

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Binary quadratic form : Class

Definition (Class)

Let ${\mathcal Q}$ be a Binary quadratic form, the following set define the class of ${\mathcal Q}$

$$\langle \mathcal{Q} \rangle = \{ S^T Q S : S \in \mathrm{Sl}(2, \mathbf{Z}) \}$$

Some notes:

- The determinant is invariant
- Every class contains a **reduced** binary quadratic form

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Reduced Binary quadratic form

Definition (Reduced form)

A binary quadratic form [A, B, C] is reduced if

 $\begin{cases} |B| \le A \le C \\ B \ge 0 \text{ if } |B| = A \text{ or } C \end{cases}$

when [A, B, C] is positive definite

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Quadratic form: Link with lattices

	Lattice formalism Quadratic form formalism	
Object	Basis $M = (u, v)$	Gram matrix $G = M^t M$
Step operation	$\textit{M} \leftarrow \textit{MS}_{\lambda}$	$G \leftarrow S_\lambda^t GS_\lambda$
Reduceness condition	$ v ^2 \ge u ^2 \ge 2 \langle u, v \rangle $	$C \ge A \ge B $

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Preliminaries

- Quantum Circuits
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Gauss Reduction

Algorithm 1 Gauss reduction

- 1: Compute Δ
- 2: while Q is not reduced do
- 3: if $|C| \leq \sqrt{|\Delta|}$ then
- 4 $t \leftarrow -sgn(C) \cdot \left| \frac{B}{2|C|} \right|$
- 5: **else**

6:
$$t \leftarrow -sgn(C) \cdot \left\lfloor \frac{\sqrt{|\Delta|+B}}{2|C|} \right\rfloor$$

7: end if
$$(0, 1)$$

8:
$$S \leftarrow \begin{pmatrix} 0 & 1 \\ 1 & t \end{pmatrix}$$

9:
$$Q \leftarrow S' QS$$

10: end while

11: return Q

Quantum Quadratic form reduction $\circ \circ \bullet \circ$

Quantum multiplication by 2^x

From Multiplication to 2^x operations

Algorithm 2 Positive Definite Reduction 1: $m \leftarrow 0$, $\epsilon \leftarrow \operatorname{sgn}(B)$ 2: if C < A then 3: $(C, A) \leftarrow (A, C)$ 4: end if 5: if $\neg (|B| < 2A)$ then 6: $m \leftarrow 2\lfloor \log_2 |B| \rfloor - \lfloor \log_2 A \rfloor - 1$ 7: else 8: $m \leftarrow \left| \frac{|B|}{2A} \right|$ 9: end if 10: if m = 0 then return $[A, (-1)^{\delta(A=-B)}B, C]$ 11: 12: else Reduce($C - \epsilon mB + m^2A, B - \epsilon 2mA, A$) 13: 14: end if

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Quantum multiplication by 2^{\times}

More global optimisations

Some more optimisations can be done:

- Iterative version : independance from the input
- Fewer conditions branches : reduce the quantum cost

We now focus on implementing the following subroutines:

- 2^x multiplication
- (Integer logarithm)

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Quantum multiplication by 2^x ••••••••

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Quantum Quadratic form reduction

Aimed Circuit



Figure: The quantum circuit for bit rotation.

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Simplified circuit

$$\begin{array}{c} |A\rangle & \hline \\ cS_m & |A \ll (2^m \cdot b)\rangle \\ |b\rangle & \hline \\ |b\rangle \end{array}$$

Figure: The quantum circuit cS_m ($A \in \{0, 1\}^n$ and $b \in \{0, 1\}$).

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Even more simplified circuit

$$|A\rangle - S_m - |A \ll 2^m\rangle$$

Figure: The quantum circuit S_m ($A \in \{0, 1\}^n$).

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Quantum multiplication by 2^x 00000000000

Example of 1-bit rotation



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Quantum multiplication by 2^x

Example of Constant time 1-(qu)bit rotation



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Constant time i-(qu)bit rotation

efficient i-(qu)bit rotation

The cyclic permutation on n qubits of parameter 2^i can be implemented :

- Depth : O(1)
- Ancilae: none
- Volume : O(*n*)

$$\sigma_0 = \left((1,0) \prod_{i=1}^{n/2-1} (i+1, n-i) \right) \cdot \left(\prod_{i=0}^{n/2-2} (i+1, n-i-1) \right)$$

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From circuit to controlled circuit

Controlled Version

Building a controlled version of S_j : cS_j can be done in with

- Depth : $O(\log n)$
- Ancilae : O(*n*)
- Volume : O(n)

 $^{0}\mbox{Cristopher}$ Moore and Martin Nilsson "Parallel quantum computation and quantum codes"

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Overall circuit for multiplication by 2^{x}



- $i = \sum_{i} i_{j} 2^{j}$ • $A \ll i = A \ll i_{0} 2^{0} \ll i_{1} 2^{1} \ll \cdots \ll i_{\log_{2}(n-1)} 2^{\log_{2}(n-1)}$
- The j-th bit of the decomposition of i in base 2 : *i_j* acts as the control bit of the circuit *cS_j*.

Complexity

Quantum Quadratic form reduction

Fine grain complexity analysis

The multiplication by a power of 2 can be inplemented on a circuit with:

- Depth : $12\lceil \log n \rceil$
- Ancilae: $n \lceil \log n \rceil$
- Volume : $12n \lceil \log n \rceil$

Results

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Algorithm	Toffoli depth	‡ Ancilae	# Toffoli gates
Modular Addition	$O(\log n)$	$O\left(\frac{n}{\log n}\right)$	O (<i>n</i>)
Multiplication	$O(n^{1.143})$	$O(n^{1.404})$	O $(n^{1.55})$
GCD	$O(n \log n)$	Ò(n)	$O(n^2)$
Bit rotation	12 log <i>n</i>	n log n	12 <i>n</i> log <i>n</i>
Logarithm	4 log <i>n</i>	4 <i>n</i>	4 <i>n</i>
Binary quadratic form Reduction	$568n\log n + 896n$	$7n^2 + 26n$	$144n^2 \log n + 2834n^2$

Conclusion

We designed the first quantum circuit that

- generalizes GCD and is very important for number theory
- performs the core step of LLL Algorithm

We also derived a fine grain complexity analysis allowing a fine estimation for applications.

Article on eprint : https://eprint.iacr.org/2022/466.pdf

Conclusion

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Thank you for your attention

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