## Quantum binary quadratic form reduction

A $\boldsymbol{O}(n \log n)$ depth circuit and application to lattice reduction

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## Human Context



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## Lattice



## Scientific Context

Lattice

- Lattice based cryptography is based on the difficulty of finding a reduced basis
- LLL is an algorithm that computes a reduced basis

Quantum Cryptanalysis

- With the help of quantum computer it is possible to break some cryptographic primitives


## How it all started



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## Efficient reversible shallow circuits

Designing efficient shallow circuits is a new domain of research. It mainly consists in

- Designing efficient circuit for basic operations
- Give its precise complexity allowing to derive precise costs in applications


## State of the art

Some previous work

- Addition
- Multiplication
- GCD

Today we will look into Binary quadratic form reduction

- first circuit designed for quadratic form reduction
- Application to LLL

[^0]
## Outline

## (1) Preliminaries

- Quantum Circuits
- Binary quadratic form


## (2) Quantum Quadratic form reduction

(3) Quantum multiplication by $2^{x}$

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## Quantum circuit

## Qubit

A qubit is the basic unit of quantum information

- pure states: $|0\rangle,|1\rangle$ (extended to $|00\rangle,|01\rangle,|11\rangle$ etc.)
- superposition : $|\psi\rangle=\alpha|0\rangle+\beta|1\rangle$ with $\alpha, \beta \in \mathbf{C}$ and $|\alpha|^{2}+|\beta|^{2}=1$


## Some Quantum Gates

- NOT.

$$
|x\rangle-\mathbb{( 1 )} \quad|\bar{x}\rangle
$$

- Bitwise addition (CNOT).

- Toffoli (CCNOT).



## Some Quantum Gates

- Swap.



## Quantum circuit and complexity measures


$|0\rangle^{k}$ are called Ancilae qubits

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$|0\rangle^{k}$ are called Ancilae qubits

- Depth: Related to the speed of the execution.
- Width: Related to the memory ( $\#$ Ancilae).
- Volume: Total number of operations.


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## Binary quadratic form : Definition

## Definition

An (integral) Binary quadratic form $\mathcal{Q}=[A, B, C]$ is a polynomial $\left(A X^{2}+B X Y+C Y^{2}\right) \in[X, Y]$. The integer $\Delta=B^{2}-4 A C$ is called the discriminant of $\mathcal{Q}$. The form is said to be:

- Degenerate when $\Delta=0$
- Positive (resp. Negative) Definite when $\Delta<0$ and

$$
\mathcal{Q}(x, y) \geq 0(\text { resp. } \mathcal{Q}(x, y) \leq 0) \text { for any }(x, y) \in \mathbf{R}^{2}
$$

Associated matrix: $Q=\left(\begin{array}{cc}A & B / 2 \\ B / 2 & C\end{array}\right)$

## Binary quadratic form : Class

## Definition (Class)

Let $\mathcal{Q}$ be a Binary quadratic form, the following set define the class of $\mathcal{Q}$

$$
\langle\mathcal{Q}\rangle=\left\{S^{\top} Q S: S \in \operatorname{Sl}(2, \mathbf{Z})\right\}
$$

Some notes:

- The determinant is invariant
- Every class contains a reduced binary quadratic form


## Reduced Binary quadratic form

## Definition (Reduced form)

A binary quadratic form $[A, B, C]$ is reduced if

$$
\left\{\begin{array}{c}
|B| \leq A \leq C \\
B \geq 0 \text { if }|B|=A \text { or } C
\end{array}\right\}
$$

when $[A, B, C]$ is positive definite

## Quadratic form: Link with lattices

|  | Lattice formalism | Quadratic form formalism |
| :---: | :---: | :---: |
| Object | Basis $M=(u, v)$ | Gram matrix $G=M^{t} M$ |
| Step operation | $M \leftarrow M S_{\lambda}$ | $G \leftarrow S_{\lambda}^{t} G S_{\lambda}$ |
| Reduceness condition | $\\|v\\|^{2} \geq\\|u\\|^{2} \geq 2\|\langle u, v\rangle\|$ | $C \geq A \geq\|B\|$ |

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## Gauss Reduction

```
Algorithm 1 Gauss reduction
    1: Compute \(\Delta\)
    while \(Q\) is not reduced do
    3: \(\quad\) if \(|C| \leq \sqrt{|\Delta|}\) then
    4: \(\quad t \leftarrow-\operatorname{sgn}(C) \cdot\left\lfloor\frac{B}{2|C|}\right\rceil\)
    5: else
    6: \(\quad t \leftarrow-\operatorname{sgn}(C) \cdot\left\lfloor\frac{\sqrt{|\Delta|}+B}{2|C|}\right\rceil\)
    7: end if
    8: \(\quad S \leftarrow\left(\begin{array}{ll}0 & 1 \\ 1 & t\end{array}\right)\)
    9: \(\quad Q \leftarrow S^{T} Q S\)
10: end while
11: return \(Q\)
```


## From Multiplication to $2^{x}$ operations

Algorithm 2 Positive Definite Reduction
1: $m \leftarrow 0, \epsilon \leftarrow \operatorname{sgn}(B)$
2: if $C<A$ then
3: $\quad(C, A) \leftarrow(A, C)$
4: end if
5: if $\neg(|B| \leq 2 A)$ then
6: $\quad m \leftarrow 2^{\left\lfloor\log _{2}|B|\right\rfloor-\left\lfloor\log _{2} A\right\rfloor-1}$
7: else
8: $\quad m \leftarrow\left\lfloor\frac{|B|}{2 A}\right\rceil$
9: end if
10: if $m=0$ then
11: return $\left[A,(-1)^{\delta(A=-B)} B, C\right]$
12: else
13: $\quad$ Reduce $\left(C-\epsilon m B+m^{2} A, B-\epsilon 2 m A, A\right)$

## 14: end if

## More global optimisations

Some more optimisations can be done:

- Iterative version : independance from the input
- Fewer conditions branches : reduce the quantum cost

We now focus on implementing the following subroutines:

- $2^{x}$ multiplication
- (Integer logarithm)


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## Aimed Circuit



Figure: The quantum circuit for bit rotation.

## Simplified circuit



Figure: The quantum circuit $c S_{m}\left(A \in\{0,1\}^{n}\right.$ and $\left.b \in\{0,1\}\right)$.

## Even more simplified circuit



Figure: The quantum circuit $S_{m}\left(A \in\{0,1\}^{n}\right)$.

## Example of 1-bit rotation



## Example of Constant time 1-(qu)bit rotation



## Constant time i-(qu)bit rotation

## efficient i-(qu)bit rotation

The cyclic permutation on $n$ qubits of parameter $2^{i}$ can be implemented :

- Depth: O (1)
- Ancilae: none
- Volume: O (n)

$$
\sigma_{0}=\left((1,0) \prod_{i=1}^{n / 2-1}(i+1, n-i)\right) \cdot\left(\prod_{i=0}^{n / 2-2}(i+1, n-i-1)\right)
$$

## From circuit to controlled circuit

## Controlled Version

Building a controlled version of $S_{j}: c S_{j}$ can be done in with

- Depth: O $(\log n)$
- Ancilae: O (n)
- Volume : O (n)

[^1]
## Overall circuit for multiplication by $2^{x}$



- $i=\sum_{i} i_{j} 2^{j}$
- $A \lll i=A \lll i_{0} 2^{0} \lll i_{1} 2^{1} \lll \cdots<i_{\log _{2}(n-1)} 2^{\log _{2}(n-1)}$
- The $j$-th bit of the decomposition of $i$ in base $2: i_{j}$ acts as the control bit of the circuit $c S_{j}$.


## Complexity

## Fine grain complexity analysis

The multiplication by a power of 2 can be inplemented on a circuit with:

- Depth : $12\lceil\log n\rceil$
- Ancilae: $n\lceil\log n\rceil$
- Volume : $12 n\lceil\log n\rceil$


## Results

| Algorithm | Toffoli depth | $\sharp$ Ancilae | $\sharp$ Toffoli gates |
| :---: | :---: | :---: | :---: |
| Modular Addition | $\mathrm{O}(\log n)$ | $\mathrm{O}\left(\frac{n}{\log n}\right)$ | $\mathrm{O}(n)$ |
| Multiplication | $\mathrm{O}\left(n^{1.143}\right)$ | $\mathrm{O}\left(n^{1.404}\right)$ | $\mathrm{O}\left(n^{1.55}\right)$ |
| GCD | $\mathrm{O}(n \log n)$ | $\mathrm{O}(n)$ | $\mathrm{O}\left(n^{2}\right)$ |
| Bit rotation | $12 \log n$ | $n \log n$ | $12 n \log n$ |
| Logarithm | $4 \log n$ | $4 n$ | $4 n$ |
| Binary quadratic form Reduction | $568 n \log n+896 n$ | $7 n^{2}+26 n$ | $144 n^{2} \log n+2834 n^{2}$ |

## Conclusion

We designed the first quantum circuit that

- generalizes GCD and is very important for number theory
- performs the core step of LLL Algorithm

We also derived a fine grain complexity analysis allowing a fine estimation for applications.
Article on eprint: https://eprint.iacr.org/2022/466.pdf

## Conclusion

## Thank you for your attention


[^0]:    ${ }^{0}$ Yasuhiro Takahashi and Noboru Kunihiro. "A fast quantum circuit for addition with few qubits. Quantum Information Computation" Srijit Dutta, Debjyoti Bhattacharjee, and Anupam Chattopadhyay. "Quantum circuits for toom-cook multiplication." Mehdi Saeedi and Igor L Markov. "Quantum Circuits for GCD Computation with O(nlogn) Depth and O(n) Ancillae."

[^1]:    ${ }^{0}$ Cristopher Moore and Martin Nilsson "Parallel quantum computation and quantum codes"

